

# Proof of the principle of mathematical induction

First recall the well-ordering axiom:

**Axiom 16 (WO).**  $(S \subseteq \mathbb{N}) \wedge (S \neq \emptyset) \Rightarrow (\exists m \in S \forall x \in S \quad m \leq x)$

In the above statement of Axiom 16 we used a common convention that the exclusive disjunction  $(m < x) \oplus (m = x)$  is abbreviated as  $m \leq x$ .

In the next theorem the universe of discourse is the set  $\mathbb{Z}_+$  of positive integers.

**Theorem 1.** Let  $P(n)$  be a propositional function involving a positive integer  $n$ . Then

$$P(1) \wedge \left( \forall k (P(k) \Rightarrow P(k+1)) \right) \Rightarrow \forall n P(n)$$

*Proof.* We will prove the contrapositive:

$$\exists j \neg P(j) \Rightarrow \neg P(1) \vee \left( \exists k (P(k) \wedge \neg P(k+1)) \right) \quad (1)$$

Assume  $\exists j \neg P(j)$ . That is, assume that there exists  $j \in \mathbb{Z}_+$  such that  $\neg P(j)$ . Now consider the set

$$S = \{n \in \mathbb{Z}_+ \mid \neg P(n)\}.$$

Clearly  $S \subseteq \mathbb{Z}_+$  and  $j \in S$ . Therefore  $S \subseteq \mathbb{N}$  and  $S \neq \emptyset$ . Hence

$$(S \subseteq \mathbb{N}) \wedge (S \neq \emptyset)$$

is true. By the well-ordering axiom we conclude

$$\exists m \in S \forall x \in S \quad m \leq x \quad (2)$$

It is important to note that such  $m$  is called a *minimum* of  $S$ .

Next we make two observations about the proposition (2). First, we notice that the proposition

$$\forall x \in S \quad m \leq x$$

is equivalent to

$$\forall x \quad x \in S \Rightarrow m \leq x,$$

which is further equivalent to

$$\forall x \quad x < m \Rightarrow x \notin S.$$

Thus (2) is equivalent to

$$\exists m \in S \quad \forall x (x < m \Rightarrow x \notin S) \quad (3)$$

Second, we notice that  $m \in \mathbb{Z}_+$ . Therefore,  $(m = 1) \oplus (1 < m)$ . In other words, there are two cases for  $m$ : either  $m = 1$  or  $m > 1$ . Consider these two cases separately.

**Case 1.** Assume  $m = 1$ . Then, since  $m = 1 \in S$ , we have that  $\neg P(1)$  is true. Consequently,

$$\neg P(1) \vee \left( \exists k (P(k) \wedge \neg P(k + 1)) \right)$$

is true. Thus, we have proved the implication (1) in this case.

**Case 2.** Assume  $m > 1$ . Then  $m - 1 > 0$  and thus  $m - 1 \in \mathbb{Z}_+$ . Define  $k = m - 1$ . Then  $k \in \mathbb{Z}_+$ . Further, since  $k < m$ , (3) implies  $k \notin S$ . Since  $n \in S$  is equivalent to  $(n \in \mathbb{Z}_+) \wedge (\neg P(n))$ ,  $k \notin S$  is equivalent to  $(k \notin \mathbb{Z}_+) \vee P(k)$ . Since  $k \in \mathbb{Z}_+$ , the last disjunction implies that  $P(k)$  is true. Recall that  $k + 1 = m \in S$ . Hence  $\neg P(k + 1)$  is true. Thus we just proved that

$$\exists k (P(k) \wedge \neg P(k + 1))$$

Consequently,

$$\neg P(1) \vee \left( \exists k (P(k) \wedge \neg P(k + 1)) \right)$$

is true. Thus, we have proved the implication (1) in Case 2, as well. This completes the proof.  $\square$